

Theorem (Granville, Schinzel, F.): *An algorithm exists for determining if a given nonreciprocal polynomial $f(x) \in \mathbb{Z}[x]$ is irreducible and that runs in time $O_{r,H}(\log n (\log \log n)^2 |\log \log \log n|)$.*

- If f has no cyclotomic factor but has a reciprocal factor, then the algorithm will give an explicit reciprocal factor.
-

f does not have a cyclotomic factor

We want to compute $\gcd(f, g)$

where $g(x) = x^{\deg f} f(1/x) \neq f(x)$.

Theorem (Granville, Schinzel, F.): *There is an algorithm which takes as input two polynomials $f(x)$ and $g(x)$ in $\mathbb{Z}[x]$, each of degree $\leq n$ and height $\leq H$ and having $\leq r + 1$ nonzero terms, with at least one of $f(x)$ and $g(x)$ free of any cyclotomic factors, and outputs the value of $\gcd_{\mathbb{Z}}(f(x), g(x))$ and runs in time $O_{r,H}(\log n)$.*

$$f(x) = \sum_{j=1}^k a_j x^{d_j} \quad \rightarrow \quad F_{\mathbb{Z}}(\vec{x}) = \sum_{j=1}^k a_j x_j$$

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$$F_1(x^{d_1}, \dots, x^{d_k}) = f(x)$$

$$F_2(x^{d_1}, \dots, x^{d_k}) = g(x)$$

Lemma (Bombieri and Zannier): *Let*

$$F_1, F_2 \in \mathbb{Q}[x_1, \dots, x_k]$$

be coprime polynomials. There exists a number $c_1(F_1, F_2)$ with the following property. If $\vec{u} = \langle u_1, \dots, u_k \rangle \in \mathbb{Z}^k$, $\xi \neq 0$ is algebraic and

$$F_1(\xi^{u_1}, \dots, \xi^{u_k}) = F_2(\xi^{u_1}, \dots, \xi^{u_k}) = 0,$$

then either ξ is a root of unity or there exists a non-zero vector $\vec{v} \in \mathbb{Z}^k$ having length at most c_1 and orthogonal to \vec{u} .

$$f(x) = \sum_{j=1}^k a_j x^{d_j} \quad \rightarrow \quad F_1(\vec{x}) = \sum_{j=1}^k a_j x_j$$

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There exists a number $c_1(\vec{a}, \vec{b}, k)$ with the following property. If $f(\xi) = g(\xi) = 0$, then there exists a non-zero vector $\vec{v} \in \mathbb{Z}^k$ having length at most c_1 and orthogonal to \vec{u} .

$$\vec{u} = \langle d_1, \dots, d_k \rangle$$

Note: It is important that c_1 is computable.

Idea: The lattice of vectors orthogonal to \vec{v} is $(k-1)$ -dimensional so that there exists a vector $\langle e_1, \dots, e_{k-1} \rangle$ and a matrix \mathcal{M} in \mathbb{Z}^{k-1} satisfying

$$\begin{pmatrix} d_1 \\ d_2 \\ \vdots \\ d_k \end{pmatrix} = \mathcal{M} \cdot \begin{pmatrix} e_1 \\ e_2 \\ \vdots \\ e_{k-1} \end{pmatrix}.$$

So

$$d_i = \sum_{j=1}^{k-1} m_{ij} e_j,$$

with the $m_{ij} \in \mathbb{Z}$ bounded.

$$d_i = \sum_{j=1}^{k-1} m_{ij} e_j \quad x^{d_i} = \prod_{j=1}^{k-1} (x^{e_j})^{m_{ij}}$$

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$$F_1^{(2)}(y_1, \dots, y_{k-1}) = \sum_{i=1}^k a_i \prod_{j=1}^{k-1} y_j^{m_{ij}}$$

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$$\implies \begin{aligned} F_1^{(k)}(x^{\text{some exponent}}) &= f(x) \\ F_2^{(k)}(x^{\text{same exponent}}) &= g(x) \end{aligned}$$

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The exponents and coefficients
in $F_1^{(j)}$ and $F_2^{(j)}$ remain bounded.

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Compute $\gcd(F_1^{(k)}(x), F_2^{(k)}(x))$.

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Note we are not saying such

$F_1^{(k)}(x)$ and $F_2^{(k)}(x)$ exist.

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**One additional item for the Final Exam
(and future Comps)**

Let $f(x) \in \mathbb{R}[x]$ with $f(0) \neq 0$. Write

$$f(x) = \sum_{j=0}^n a_j x^j = a_n \prod_{j=1}^n (x - \alpha_j)$$

and

$$w(x) = a_n \prod_{\substack{1 \leq j \leq n \\ |\alpha_j| > 1}} (x - \alpha_j) \prod_{\substack{1 \leq j \leq n \\ |\alpha_j| \leq 1}} (\alpha_j x - 1).$$

Recall that

$$M(f) = |a_n| \prod_{1 \leq j \leq n} \max\{|\alpha_j|, 1\} \quad \text{and} \quad \|f\| = \sqrt{\sum_{j=0}^n a_j^2}.$$

Prove the following:

- (a) Explain why $w(x)\tilde{w}(x) = f(x)\tilde{f}(x)$.
- (b) Prove $M(f) \leq \|f\|$.
- (c) For $f(x) \in \mathbb{R}[x]$, prove $\|f\| \leq 2^{\deg f} M(f)$.
- (d) Let $f(x)$ and $g(x)$ be polynomials in $\mathbb{Z}[x]$ such that $g(x) | f(x)$.
Prove $\|g\| \leq 2^{\deg g} \|f\|$.