The Lattice Base Reduction Algorithm

This is a method which was developed in 1982 by Arjen Lenstra, Hendrik Lenstra and László Lovász to prove that factoring polynomials in $\mathbb{Z}[x]$ can be done in polynomial time. It is sometimes called the LLL-algorithm or the L³-algorithm.

Definitions and Notations. Let \mathbb{Q}^n denote the set of vectors $\langle a_1, a_2, \ldots, a_n \rangle$ with $a_j \in \mathbb{Q}$. For $\vec{b} = \langle a_1, a_2, \ldots, a_n \rangle \in \mathbb{Q}^n$ and $\vec{b}' = \langle a'_1, a'_2, \ldots, a'_n \rangle \in \mathbb{Q}^n$, define the usual dot product $\vec{b} \cdot \vec{b}'$ by

$$ec{b}\cdotec{b}'=a_1a_1'+a_2a_2'+\cdots+a_na_n',$$

and set

$$\|ec{b}\| = \sqrt{a_1^2 + a_2^2 + \dots + a_n^2}.$$

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Further, we use A^T to denote the transpose of a matrix A, so the rows and columns of A are the same as the columns and rows of A^T , respectively. Let $\vec{b}_1, \ldots, \vec{b}_n \in \mathbb{Q}^n$, and let $A = (\vec{b}_1, \ldots, \vec{b}_n)$ be the $n \times n$ matrix with column vectors $\vec{b}_1, \ldots, \vec{b}_n$. The lattice \mathcal{L} generated by $\vec{b}_1, \ldots, \vec{b}_n$ is

$$\mathcal{L}=\mathcal{L}(A)=ec{b}_1\mathbb{Z}+\dots+ec{b}_n\mathbb{Z}.$$

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Comment: Different A can determine the same \mathcal{L} . But given \mathcal{L} , the value of $|\det A|$ is the same for all such A. To see this, observe that if $\vec{b}_1, \ldots, \vec{b}_n$ and $\vec{b}'_1, \ldots, \vec{b}'_n$ are two bases for \mathcal{L} , there are matrices U and V with integer entries such that

$$ig(ec{b}_1,\ldots,ec{b}_nig)UV=ig(ec{b}_1',\ldots,ec{b}_n'ig)V=ig(ec{b}_1,\ldots,ec{b}_nig).$$

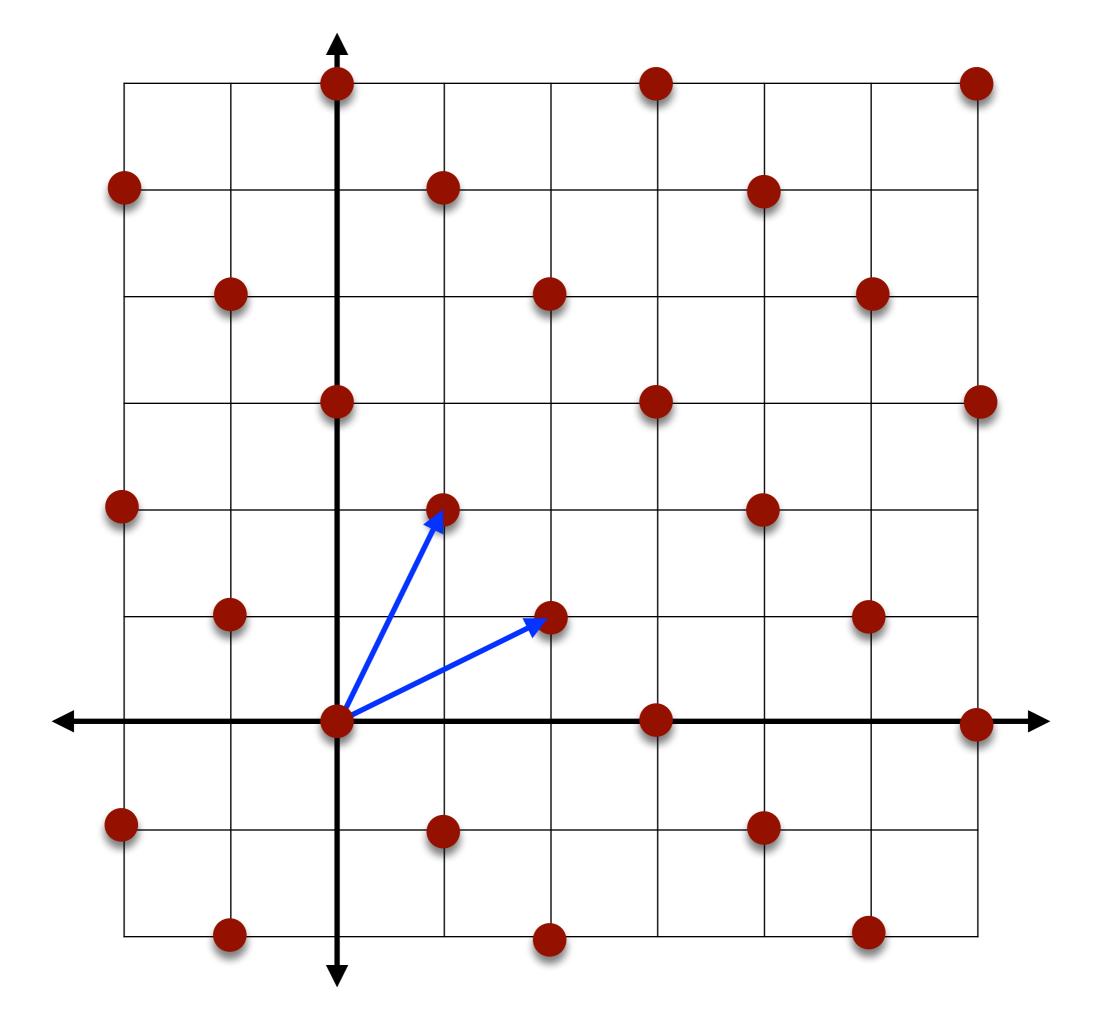
Given that $\vec{b}_1, \ldots, \vec{b}_n$ is a basis for \mathbb{R}^n , it follows that UV is the identity matrix and det $V = \pm 1$. The second equation above then implies

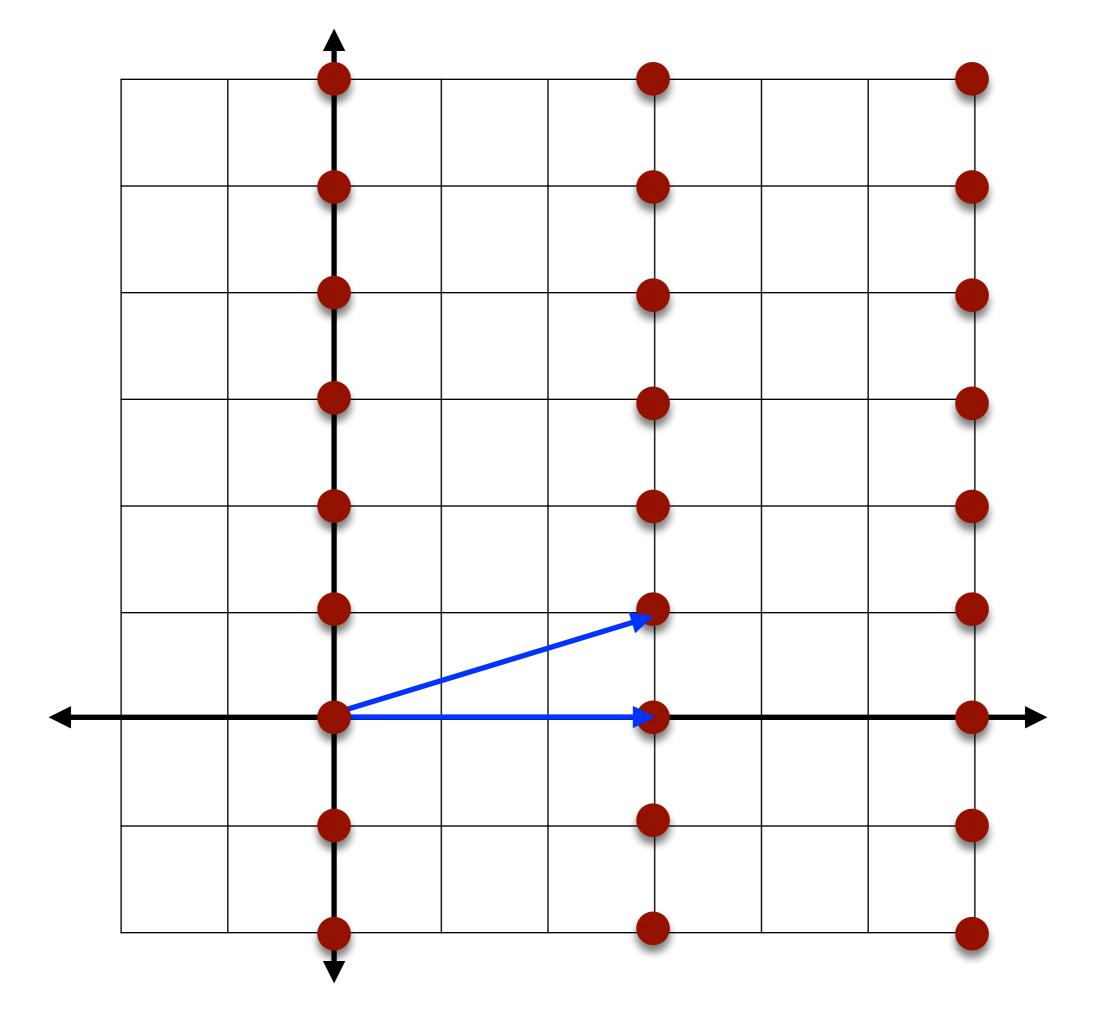
$$|\det\left(ec{b}_1',\ldots,ec{b}_n'
ight)| = |\det\left(ec{b}_1,\ldots,ec{b}_n
ight)|.$$

We set det \mathcal{L} to be this common value.

Example. In \mathbb{R}^2 , the lattice formed from the basis $\langle 1, 0 \rangle$ and $\langle 0, 1 \rangle$ is the same as the lattice formed from the basis $\langle 1, 0 \rangle$ and $\langle 1, 1 \rangle$. This can be seen geometrically and algebraically.

Example 2. The lattice \mathcal{L}_1 with basis $\langle 2, 1 \rangle$ and $\langle 1, 2 \rangle$ and the lattice \mathcal{L}_2 with basis $\langle 3, 0 \rangle$ and $\langle 3, 1 \rangle$ are such that det $\mathcal{L}_1 = \det \mathcal{L}_2$. But the lattices are quite different.





The Gram-Schmidt orthogonalization process Define recursively

$$ec{b}_i^* = ec{b}_i - \sum_{j=1}^{i-1} \mu_{ij} ec{b}_j^*, \qquad ext{for } 1 \leq i \leq n,$$

where

$$\mu_{ij} = \mu_{i,j} = rac{ec{b}_i \cdot ec{b}_j^*}{ec{b}_j^* \cdot ec{b}_j^*}, \qquad ext{for } 1 \leq j < i \leq n.$$

Then for each $i \in \{1, \ldots, n\}$, the vectors $\vec{b}_1^*, \ldots, \vec{b}_i^*$ span the same subspace of \mathbb{R}^n as $\vec{b}_1, \ldots, \vec{b}_i$. In other words,

$$egin{aligned} &\{a_1ec{b}_1^*+\dots+a_iec{b}_i^*:a_j\in\mathbb{R} ext{ for }1\leq j\leq i\}\ &=\{a_1ec{b}_1+\dots+a_iec{b}_i:a_j\in\mathbb{R} ext{ for }1\leq j\leq i\}. \end{aligned}$$

Furthermore, the vectors $\vec{b}_1^*, \ldots, \vec{b}_n^*$ are linearly independent (hence, non-zero) and pairwise orthogonal (i.e., for distinct i and j, we have $\vec{b}_i^* \cdot \vec{b}_j^* = 0$).

Hadamard's Inequality

The value of det \mathcal{L} can be viewed as the volume of the polyhedron with edges parallel to and the same length as $\vec{b}_1, \ldots, \vec{b}_n$. This volume is independent of the basis that is used for \mathcal{L} . Geometrically (in low dimensions),

 $\det \mathcal{L} \leq \|ec{b}_1\| \, \|ec{b}_2\| \cdots \|ec{b}_n\|.$

$$\begin{split} \vec{b} \in \mathcal{L}, \ \vec{b} \neq 0 \implies \|\vec{b}\| \ge \min\{\|\vec{b}_1^*\|, \|\vec{b}_2^*\|, \dots, \|\vec{b}_n^*\|\} \\ \vec{b}_i = \vec{b}_i^* + \sum_{j=1}^{i-1} \mu_{ij} \vec{b}_j^* \qquad \mu_{i,j} = \frac{\vec{b}_i \cdot \vec{b}_j^*}{\vec{b}_j^* \cdot \vec{b}_j^*} \\ \vec{b} = u_1 \vec{b}_1 + \dots + u_k \vec{b}_k, \quad \text{where each } u_j \in \mathbb{Z} \text{ and } u_k \neq 0 \\ \vec{b} = v_1 \vec{b}_1^* + \dots + v_k \vec{b}_k^*, \quad \text{where each } v_j \in \mathbb{Q} \text{ and } v_k = u_k \\ \|\vec{b}\|^2 = (v_1 \vec{b}_1^* + \dots + v_k \vec{b}_k^*) \cdot (v_1 \vec{b}_1^* + \dots + v_k \vec{b}_k^*) \\ = v_1^2 \|\vec{b}_1^*\|^2 + \dots + v_k^2 \|\vec{b}_k^*\|^2 \ge \|\vec{b}_k^*\|^2 \\ \vec{b} \in \mathcal{L}, \ k \text{ as above } \implies \|\vec{b}\|^2 \ge \|\vec{b}_k^*\|^2 \end{split}$$

(*)

Definition. Let $\vec{b}_1, \ldots, \vec{b}_n$ be a basis for a lattice \mathcal{L} , and let $\vec{b}_1^*, \ldots, \vec{b}_n^*$ be the corresponding basis in \mathbb{R}^n obtained from the Gram-Schmidt orthogonalization process, with μ_{ij} as defined before. Then we say that $\vec{b}_1, \ldots, \vec{b}_n$ is *reduced* if both of the following hold

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$$\begin{array}{ll} \text{(i)} \ |\mu_{ij}| \leq \frac{1}{2} & \text{for } 1 \leq j < i \leq n \\ \\ \text{(ii)} \ \|\vec{b}_i^* + \mu_{i,i-1}\vec{b}_{i-1}^*\|^2 \geq \frac{3}{4} \, \|\vec{b}_{i-1}^*\|^2 & \text{for } 1 < i \leq n. \end{array}$$

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$$ec{b}\in\mathcal{L}, \ ec{b}
eq 0 \ \Longrightarrow \ \|ec{b}_1\|\leq 2^{(n-1)/2}\|ec{b}\|.$$

Thus, \vec{b}_1 is not far from being the shortest vector in \mathcal{L} .

$$\|ec{b}_{i}^{*}\|^{2} + rac{1}{4}\|ec{b}_{i-1}^{*}\|^{2} \geq \|ec{b}_{i}^{*} + \mu_{i,i-1}ec{b}_{i-1}^{*}\|^{2} \geq rac{3}{4}\|ec{b}_{i-1}^{*}\|^{2}$$

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(ii) $\|\vec{b}_i^* + \mu_{i,i-1}\vec{b}_{i-1}^*\|^2 \geq \frac{3}{4}\|\vec{b}_{i-1}^*\|^2$ for $1 < i \leq n$.

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$$\|\vec{b}_i^*\|^2 + \frac{1}{4} \|\vec{b}_{i-1}^*\|^2 \geq \|\vec{b}_i^* + \mu_{i,i-1}\vec{b}_{i-1}^*\|^2 \geq \frac{3}{4} \|\vec{b}_{i-1}^*\|^2 \implies \|\vec{b}_i^*\|^2 \geq (1/2) \|\vec{b}_{i-1}^*\|^2 \\ \implies \|\vec{b}_i^*\|^2 \geq \frac{1}{2^{i-j}} \|\vec{b}_j^*\|^2 \quad \text{for } 1 \leq j < i \leq n \\ \|\vec{b}\|^2 \geq \|\vec{b}_k^*\|^2 \geq \frac{1}{2^{k-1}} \|\vec{b}_1^*\|^2 \geq \frac{1}{2^{n-1}} \|\vec{b}_1^*\|^2 = \frac{1}{2^{n-1}} \|\vec{b}_1\|^2 \end{split}$$

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Let $\vec{x}_1, \vec{x}_2, \dots, \vec{x}_t$ be t linearly independent vectors in \mathcal{L} . Then $\|\vec{b}_j\| \leq 2^{(n-1)/2} \max\{\|\vec{x}_1\|, \|\vec{x}_2\|, \dots, \|\vec{x}_t\|\}$ for $1 \leq j \leq t$.

$$\|ec{b}_i^*\|^2 \geq rac{1}{2^{i-j}}\|ec{b}_j^*\|^2 \quad ext{ for } 1 \leq j < i \leq n$$

$$\|ec{b}_i\|^2 = \left\|ec{b}_i^* + \sum_{j=1}^{i-1} \mu_{ij}ec{b}_j^*
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 $\|ec{b}_{j}\|^{2} \leq 2^{j-1} \|ec{b}_{j}^{*}\|^{2} \leq 2^{i-1} \|ec{b}_{i}^{*}\|^{2} \quad ext{ for } 1 \leq j \leq i \leq n$

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$$\begin{split} \|\vec{b}_{j}\|^{2} &\leq 2^{j-1} \|\vec{b}_{j}^{*}\|^{2} \leq 2^{i-1} \|\vec{b}_{i}^{*}\|^{2} \quad \text{ for } 1 \leq j \leq i \leq n \\ (*) \qquad \vec{b} \in \mathcal{L}, \ k \text{ as above } \implies \|\vec{b}\|^{2} \geq \|\vec{b}_{k}^{*}\|^{2} \\ \vec{x}_{j} &= \sum_{i=1}^{m(j)} u_{ji}\vec{b}_{i}, \quad u_{jm(j)} \neq 0, \quad m(1) \leq m(2) \leq \cdots \leq m(t) \\ \qquad m(j) \geq j \text{ for } 1 \leq j \leq t \\ \|\vec{x}_{j}\|^{2} \geq \|\vec{b}_{m(j)}^{*}\|^{2} \quad \text{ for } 1 \leq j \leq t \\ \|\vec{b}_{j}\|^{2} \leq 2^{m(j)-1} \|\vec{b}_{m(j)}^{*}\|^{2} \leq 2^{n-1} \|\vec{x}_{j}\|^{2} \quad \text{ for } 1 \leq j \leq t \end{split}$$

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Thus, from Hadamard's inequality, we obtain

 $2^{-n(n-1)/4} \|ec{b}_1\| \, \|ec{b}_2\| \cdots \|ec{b}_n\| \leq \det \mathcal{L} \leq \|ec{b}_1'\| \, \|ec{b}_2'\| \cdots \|ec{b}_n'\|$ for any basis $ec{b}_1', \dots, ec{b}_n'$ of \mathcal{L} .

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Comment: Recall that finding a basis $\vec{b}'_1, \ldots, \vec{b}'_n$ for which the product on the right is minimal is NP-hard. The above implies that a reduced basis is close to being such a basis.

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New Goal: Show how one can determine $h_0(x)$ using h(x)and without worrying about other factors of f(x) modulo p^k .

Why would this improve on the Zassenhaus approach?

What is the lattice we want to use?

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 $h(x) ext{ monic irreducible factor of } f(x) ext{ modulo } p^k \ h_0(x) | f(x) ext{ in } \mathbb{Z}[x], ext{ } h(x) | h_0(x) ext{ modulo } p^k$

$$\ell = \deg h, \; m \in \{\ell, \ell+1, \ldots, n-1\}$$

 $m ext{ is the possible degree of } h_0(x)$

$$egin{aligned} w(x) &= a_m x^m + \dots + a_1 x + a_0 \in \mathbb{Z}[x] \ & \longleftrightarrow \quad ec{b} &= \langle a_0, a_1, \dots, a_m
angle \in \mathbb{Z}^{m+1} \end{aligned}$$

Define \mathcal{L} to be the lattice in \mathbb{Z}^{m+1} spanned by the vectors associated with

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Example
$$\begin{bmatrix} > f := x^{14}-4*x^{3}+2*x^{2}+x-3; \\ f := x^{14}-4x^{3}+2x^{2}+x-3 \end{bmatrix}$$
$$\begin{bmatrix} > Factor(f) \mod 151; \\ (x^{2}+129x+44) (x^{2}+147x+92) (x^{2}+127x+31) (x^{7}+24x^{6}+91x^{5}+81x^{4}+30x^{3}+20x^{2}+2x+34) (x+26) \end{bmatrix}$$

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$$(x^{2} + 129 x + 44) (x^{2} + 147 x + 92) (x^{2} + 127 x + 31) (x^{7} + 24 x^{6} + 91 x^{5} + 81 x^{4} + 30 x^{3} + 20 x^{2} + 2 x + 34) (x + 26)$$

 $\langle 151, 0, 0, 0, 0, 0 \rangle$ $\langle 0, 151, 0, 0, 0, 0 \rangle$ $\langle 44, 129, 1, 0, 0, 0 \rangle$ $\langle 0, 44, 129, 1, 0, 0 \rangle$ $\langle 0, 0, 44, 129, 1, 0 \rangle$ $\langle 0, 0, 0, 44, 129, 1 \rangle$

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Claim: The lattice \mathcal{L} is exactly the vectors corresponding to $w(x) \in \mathbb{Z}[x]$ that are divisible by h(x) modulo p^k . $\langle 151, 0, 0, 0, 0, 0 \rangle$ $\langle 0, 151, 0, 0, 0, 0 \rangle$ $\langle 44, 129, 1, 0, 0, 0 \rangle$ $\langle 0, 44, 129, 1, 0, 0 \rangle$

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$$b_1 = \langle a_0, a_1, \dots, a_m
angle,$$

then

$$ec{b}_0 = \langle a_0/d, a_1/d, \dots, a_m/d
angle,$$
where $d = \gcd(a_0, \dots, a_m).$

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\langle 0, 0, 0, 44, 129, 1 \rangle
```

Claim: The lattice \mathcal{L} is exactly the vectors corresponding to $w(x) \in \mathbb{Z}[x]$ that are divisible by h(x) modulo p^k . Hence, the vector \vec{b}_0 corresponding to $h_0(x)$ is in \mathcal{L} .

(Go to Maple.)

> f :=
$$x^{14} - 4x^{3} + 2x^{2} + x - 3$$
;
f := $x^{14} - 4x^3 + 2x^2 + x - 3$
> Factor(f) mod 151;
($x^2 + 129x + 44$) ($x^2 + 147x + 92$) (x^2
+ 127x + 31) ($x^7 + 24x^6 + 91x^5 + 81x^4$
+ 30 $x^3 + 20x^2 + 2x + 34$) (x + 26)

We will show that in fact if p^k is large and $\vec{b}_1, \ldots, \vec{b}_{m+1}$ is a reduced basis for \mathcal{L} with

$$b_1 = \langle a_0, a_1, \dots, a_m
angle,$$

then

$$ec{b}_0 = \langle a_0/d, a_1/d, \dots, a_m/d
angle,$$
where $d = \gcd(a_0, \dots, a_m).$