## $\S$ **2.2 Eisenstein Polynomials**

We say that a polynomial  $f(x) = \sum_{j=0}^{n} a_j x^j \in \mathbb{Z}[x]$  is in Eisenstein form (with respect to the prime p) if there is a prime p such that  $p \nmid a_n$ ,  $p|a_j$  for j < n, and  $p^2 \nmid a_0$ . An Eisenstein polynomial is an  $f(x) \in \mathbb{Z}[x]$  for which there is an integer a and a prime p such that f(x + a) is in Eisenstein form with respect to the prime p. In other words,  $f(x) \in \mathbb{Z}[x]$  is Eisenstein if there is an integer a and a prime p such that  $f(x + a) = \sum_{j=0}^{n} a'_j x^j$  where  $p \nmid a'_n$ ,  $p|a'_j$  for j < n, and  $p^2 \nmid a'_0$ . More specifically, we say that such an f(x) is Eisenstein with respect to the prime p. For example, since  $f(x) = x^2 + x + 1$  is such that  $f(x + 1) = x^2 + 3x + 3$ , the polynomial f(x) is Eisenstein with respect to 3. It follows easily from Theorem 2.1.1 that if f(x) is Eisenstein with respect to a prime p, then f(x) is irreducible over  $\mathbb{Q}$  (see Exercise (1.1)).

Suppose one is given an  $f(x) \in \mathbb{Z}[x]$  and wishes to decide whether f(x) is Eisenstein with respect to some prime (which is not given). We assume  $n = \deg f(x)$  is at least 2. One approach to making such a decision involves the use of discriminants or resultants. Our presentation here will be restricted to resultants. Let  $f(x) = \sum_{j=0}^{n} a_j x^j$  and  $g(x) = \sum_{j=0}^{r} b_j x^j$  be in  $\mathbb{C}[x]$  with  $n \ge 1, r \ge 1$  and  $a_n b_r \ne 0$ . We define the resultant of f(x) and g(x) in terms of the Sylvester determinant R(f,g) associated with f(x) and g(x). The resultant R(f,g) is the determinant of an  $(n+r) \times (n+r)$  matrix with the first r rows consisting of the coefficients of f(x), where each of these rows contains one more leading 0 than its predecessor, and with the last n rows consisting of the coefficiently, we have <sup>1</sup>

$$(2.2.1) R(f,g) = \begin{vmatrix} a_n & a_{n-1} & a_{n-2} & \dots & a_0 & 0 & 0 & \dots & 0 \\ 0 & a_n & a_{n-1} & \dots & a_1 & a_0 & 0 & \dots & 0 \\ 0 & 0 & a_n & \dots & a_2 & a_1 & a_0 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \ddots & \vdots \\ b_r & b_{r-1} & b_{r-2} & \dots & b_0 & 0 & 0 & \dots & 0 \\ 0 & b_r & b_{r-1} & \dots & b_1 & b_0 & 0 & \dots & 0 \\ 0 & 0 & b_r & \dots & b_2 & b_1 & b_0 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \ddots & \vdots \end{vmatrix}$$

For example, if  $f(x) = x^3 + x + 1$  and  $g(x) = 2x^2 + x + 3$ , then (2.2.1) becomes

$$R(f,g) = \begin{vmatrix} 1 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 2 & 1 & 3 & 0 & 0 \\ 0 & 2 & 1 & 3 & 0 \\ 0 & 0 & 2 & 1 & 3 \end{vmatrix}.$$

<sup>&</sup>lt;sup>1</sup>The appearance of the right-hand side of (2.2.1) is somewhat misleading. The entry  $a_0$  in the first row, for example, is not necessarily in the same column as the entry  $b_0$  in the first row consisting of the  $b_i$ 's.

**Lemma 2.2.1.** Let f(x) and  $g(x) \in \mathbb{C}[x]$ , and suppose that there is an  $\alpha$  such that  $f(\alpha) = g(\alpha) = 0$ . Then R(f,g) = 0.

**Proof.** Add to the *i*th row of the last column (the (n + r)th column) of the determinant on the right-hand side of (2.2.1) the product of the entry in the *i*th row and *j*th column with  $\alpha^{n+r-j}$ . Then the first *r* entries in the last column become  $\alpha^{r-1}f(\alpha), \alpha^{r-2}f(\alpha), \ldots, f(\alpha)$  and the last *n* entries become  $\alpha^{n-1}g(\alpha), \alpha^{n-2}g(\alpha), \ldots, g(\alpha)$ . By the conditions of the lemma, these are all 0, and the result follows.

Before continuing, it is of interest to point out that the above proof can be modified slightly to obtain another result of interest. We replace the role of  $\alpha$  above with a variable x, adjusting the determinant in (2.2.1) so that the right-most column consists of the entries

(2.2.2) 
$$x^{r-1}f(x), x^{r-2}f(x), \dots, f(x), x^{n-1}g(x), x^{n-2}g(x), \dots, g(x).$$

The other entries in (2.2.1) remain untouched and, hence, are coefficients of the polynomials f(x) and g(x). Call the  $(n + r) \times (n + r)$  matrix associated with this determinant A. Expanding det A along the right-most column, we obtain

(2.2.3) 
$$f(x)u(x) + g(x)v(x) = R(f,g),$$

for some polynomials u(x) and v(x) with deg  $u < \deg g$  and deg  $v < \deg f$ . Of particular significance here is that if f(x) and g(x) are in  $\mathbb{Z}[x]$ , then (2.2.3) is a linear combination of f(x) and g(x) in the ring  $\mathbb{Z}[x]$ . In other words, in this case, u(x) and v(x) in (2.2.3) are in  $\mathbb{Z}[x]$ . The problem of finding the smallest positive integer d that can be so represented as a linear combination of two given relatively prime polynomials f(x) and g(x) in  $\mathbb{Z}[x]$  is non-trivial. We will examine this in a later chapter in the special context of f(x) and g(x) being cyclotomic polynomials.

Of interest to us in (2.2.3) is the case that f(x) and g(x) are non-constant polynomials and R(f,g) = 0 over the field of rationals or over the field of arithmetic modulo a prime p. Denote the field by F. Our argument for (2.2.3) still holds over F except that we would like to know that u(x) and v(x) are not identically zero in our field. For this purpose, we observe that R(f,g) = 0corresponds to det A = 0 which corresponds to the rows of A being linearly dependent over F. Thus, if the *j*th row of A is the vector  $\overrightarrow{v}_j$  consisting of n+rcomponents, then there are  $c_j \in F$  not all zero such that

$$c_1 \overrightarrow{v}_1 + c_2 \overrightarrow{v}_2 + \dots + c_{n+r} \overrightarrow{v}_{n+r} = \overrightarrow{0}.$$

In particular, recalling that the entries in the last column of A are given by (2.2.2), we see that (2.2.3) holds with

$$u(x) = c_1 x^{r-1} + c_2 x^{r-2} + \dots + c_r$$

and

$$v(x) = c_{r+1}x^{n-1} + c_{r+2}x^{n-2} + \dots + c_{r+n}$$

Here, we have u(x) and v(x) are in F[x], at least one of u(x) and v(x) is non-zero (hence, both are), deg  $u < \deg g$  and deg  $v < \deg f$ . With such u(x) and v(x) in hand, we are now ready to show the following (but note the cautionary remark after the proof).

**Lemma 2.2.2.** Let f(x) and g(x) be two non-constant polynomials in the field F of rational numbers or arithmetic modulo a prime p. If R(f,g) is zero in F, then f(x) and g(x) have an irreducible factor in common in F[x]. If further deg  $g < \deg f$ , then f(x) is reducible over F.

**Proof.** From (2.2.3) and R(f,g) = 0, we deduce f(x)u(x) = -g(x)v(x) in F[x]. Since v(x) is non-zero with degree less than the degree of f(x), unique factorization in F[x] implies the conclusions of the lemma.

To clarify, the polynomials in Lemma 2.2.2 are to be viewed as polynomials in the field F. In particular, if one has polynomials f(x) and g(x) in  $\mathbb{Z}[x]$  and wishes to apply the lemma modulo a prime p, then the polynomials should be reduced modulo p, possibly changing the degrees of the polynomials, before computing R(f,g) with the Sylvester determinant and applying Lemma 2.2.2. By way of an example, consider

$$f(x) = 2x^3 + x^2 + x + 1$$
 and  $g(x) = 2x^2 + x + 1$ .

In this case, one can check that  $R(f,g) = 8 = 2^3$ . Observe, however, that f(x) and g(x) do not have an irreducible factor in common modulo 2. Indeed, we have

$$f(x) \equiv x^2 + x + 1 \pmod{2}, \quad g(x) \equiv x + 1 \pmod{2}$$

and

$$R(x^2 + x + 1, x + 1) = 1.$$

If  $\alpha_1, \ldots, \alpha_n$  are the roots of f(x), one can show that

(2.2.4) 
$$R(f,g) = a_n^r g(\alpha_1) \cdots g(\alpha_n).$$

The proof can be found in Uspensky (1948). Observe that (2.2.4) implies Lemma 2.2.1 and also the converse of Lemma 2.2.1. Thus, if R(f,g) = 0, then f(x) and g(x) have a complex root in common. With the cautionary notes of the previous paragraph, we can view Lemma 2.2.2 as a consequence of (2.2.4).

We now show that the following algorithm works to determine whether a given polynomial f(x) is Eisenstein.

**Algorithm:** Suppose that  $f(x) \in \mathbb{Z}[x]$  is of degree  $n \geq 2$ . Calculate R(f, f') (using the right-hand side of (2.2.1)). If R(f, f') = 0, then f(x) is not Eisenstein with respect to any prime. If  $R(f, f') \neq 0$ , then factor it. For each

prime p dividing R(f, f'), check to see if any of the translates f(x + a), where  $a \in \{0, 1, ..., p - 1\}$ , is in Eisenstein form with respect to the prime p. If such a prime p and such an a are such that f(x + a) is in Eisenstein form with respect to p, then f(x) is Eisenstein with respect to p. If no such prime p and no such a are such that f(x + a) is in Eisenstein form with respect to p, then f(x) is not Eisenstein form with respect to any prime.

In justifying the algorithm, we explain how one can use the resultant R(f, f') to determine whether a polynomial f(x) has a multiple factor (a factor which appears with multiplicity > 1) modulo some prime (which is unspecified). To see this, suppose that there is a prime p such that

(2.2.5) 
$$f(x) \equiv g(x)^2 h(x) \pmod{p}$$

where g(x) is of degree  $\geq 1$ . Note that if for some integer a we have that f(x+a) is in Eisenstein form with respect to the prime p, then  $f(x) \equiv a_n(x-a)^n \pmod{p}$  so that one can take g(x) = x - a. Define  $f_1(x) = g(x)^2 h(x)$  so that the coefficients of f(x) and of f'(x) are the same as the corresponding coefficients of  $f_1(x)$  and  $f'_1(x)$  all considered modulo p. In particular,  $R(f, f') \equiv R(f_1, f'_1) \pmod{p}$ . Since

$$f_1'(x) = 2g(x)g'(x)h(x) + g(x)^2h'(x) = g(x)\left(2g'(x)h(x) + g(x)h'(x)\right),$$

we get that each root of g(x) is a root of  $f_1(x)$  and of  $f'_1(x)$ . By Lemma 2.2.1, we get that  $R(f_1, f'_1) = 0$ . Hence,  $R(f, f') \equiv 0 \pmod{p}$ . Thus, p divides R(f, f'); and to determine if (2.2.5) holds for some prime p, we simply need to check whether it holds for each prime divisor p of R(f, f'). The fact that the algorithm works when  $R(f, f') \neq 0$  is now fairly straight forward, but we need to justify that we can restrict our consideration of integers a to  $a \in \{0, 1, \ldots, p-1\}$ . For this purpose, we suppose that b is an integer for which f(x+b) is in Eisenstein form with respect to some prime p and show that f(x+a) is also for any  $a \equiv b \pmod{p}$ . Since  $f(x+b) \equiv f(x+a) \pmod{p}$ , we simply need to justify that  $p^2$  does not divide the constant term in f(x+a). In other words, we want to show that  $p^2 \nmid f(a)$ . Writing  $f(x+b) = \sum_{j=0}^{n} a'_j x^j$ , we get that  $p|a'_j$  for j < n and  $p^2 \nmid a'_0$ . Writing a = kp + b where k is an integer, we get that

$$f(a) \equiv f(kp+b) \equiv \sum_{j=0}^{n} a'_{j} k^{j} p^{j} \equiv kpa'_{1} + a'_{0} \equiv a'_{0} \pmod{p^{2}}.$$

Thus,  $f(a) \not\equiv 0 \pmod{p^2}$ , completing what we set out to show (for the case  $R(f, f') \neq 0$ ).

If R(f, f') = 0, the above all works except that every prime is a prime divisor of R(f, f') so it is not reasonable to consider all the prime divisors of R(f, f'). But observe that (2.2.4) implies that f(x) and f'(x) have a root in common. Hence, in this case, f(x) must have a multiple root (the reader should justify this) so that f(x) is reducible (see Exercise (1.8)). Alternatively, Lemma 2.2.2 implies that f(x) is reducible over  $\mathbb{Q}$ . In particular, by Theorem 2.1.1, we can conclude that f(x) cannot be Eisenstein with respect to any prime.

**Example:** Consider  $f(x) = x^4 + 2x - 1$ , and suppose that we wish to find every prime p such that f(x) is Eisenstein with respect to p. We first calculate R(f, f') by using (2.2.1). To do this somewhat efficiently, we multiply below the first row by -4 and add it to the fourth row. Observe that we will get the same result in the fifth row with an extra leading 0 if we multiply the second row by -4 and add it to the fifth row. Similarly, we can obtain the same result in the sixth row (as the fourth row) with 2 extra leading 0's by considering the third row. We get that

$$R(f,f') = \begin{vmatrix} 1 & 0 & 0 & 2 & -1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 2 & -1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 2 & -1 \\ 4 & 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 4 & 0 & 0 & 2 & 0 & 0 \\ 0 & 0 & 4 & 0 & 0 & 2 & 0 \\ 0 & 0 & 0 & 4 & 0 & 0 & 2 \end{vmatrix} = \begin{vmatrix} 1 & 0 & 0 & 2 & -1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 2 & -1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 2 & -1 \\ 0 & 0 & 0 & -6 & 4 & 0 & 0 \\ 0 & 0 & 0 & 0 & -6 & 4 & 0 \\ 0 & 0 & 0 & 0 & 0 & -6 & 4 \\ 0 & 0 & 0 & 0 & 4 & 0 & 0 & 2 \end{vmatrix}.$$

A direct computation now gives

$$R(f, f') = \begin{vmatrix} -6 & 4 & 0 & 0\\ 0 & -6 & 4 & 0\\ 0 & 0 & -6 & 4\\ 4 & 0 & 0 & 2 \end{vmatrix} = -6 \times 72 - 4 \times 64 = -688$$

Since  $688 = 2^4 \times 43$ , we only need to deal with the primes 2 and 43. We make use of Exercise (2.6). Observe that 2 divides f(1), and so we consider  $f(x+1) = x^4 + 4x^3 + 6x^2 + 6x + 2$ . Thus, f(x) is Eisenstein with respect to the prime 2 (and, hence, f(x) is irreducible). Observe that 43 divides f(3) but that  $f'(3) = 4 \times 27 + 2 = 110$  is not divisible by 43. Thus, f(x) is not Eisenstein with respect to the prime 43. Hence, 2 is the only prime p such that f(x) is Eisenstein with respect to p. Alternatively, we note that Exercise (2.13) could have been used to determine that f(x) is not Eisenstein with respect to 43.

In this section, we have considered the problem of determining whether a polynomial  $f(x) \in \mathbb{Z}[x]$  can under a translation be shown to be irreducible by the Schönemann-Eisenstein criterion. In general, if f(x) and  $g(x) \in \mathbb{Z}[x]$  and f(g(x)) is irreducible, then f(x) is irreducible; hence, it is reasonable to attempt to determine whether a given f(x) is irreducible by applying the Schönemann-Eisenstein criterion after composing f(x) with another polynomial. We leave further consideration of this idea as an exercise (Exercise (2.10)).